# SAT/SMT solving in Haskell

### ろ Masahiro Sakai (酒井 政裕) Haskell Day 2016 2016-09-17

### Self Introduction Masahiro Sakai

- \* Twitter: @masahiro\_sakai
  \* github: https://github.com/msakai/
  \* G+: https://plus.google.com/+MasahiroSakai
- Translated "Software Abstractions" and TaPL into Japanese with colleagues
   Interests: Categorical Programming, Theorem Proving / Decision Procedures,

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> > Software Abstractio

型システム入門 プログラミング言語と型の理論

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# \* What are SAT and SMT? \* Haskell libraries for SMT solving



# \* toysat/toysmt \* Conclusion

# What are SAT and SM1?

# What is SAT?

### \* \* SAT = Boolean SATisfiability problem

- \* "Is there an assignment that makes given formula true?"
- \* Examples:
  - \*  $(P_{\vee}Q)_{\wedge}(P_{\vee} \neg Q)_{\wedge}(\neg P_{\vee} \neg Q)$  is satisfiable with
    - ${P \mapsto True, Q \mapsto False}$
  - \*  $(P \lor Q) \land (P \lor \neg Q) \land (\neg P \lor \neg Q) \land (\neg P \lor Q)$  is unsatisfiable
- SAT is NP complete, but state-of-the-art SAT-solver can often solve problems with millions of variables / constraints.

# What is SM1?

- \* Weakness of SAT: Really low-level representation
  - \* Encoding problems into SAT sometimes blows-up
  - \* SAT solver cannot leverage high-level knowledge
- \* SMT = Satisfiability Modulo Theories
  - \* An approach to overcome the weakness of SAT
  - ★ Problem Example: Is there array a, function f, integers i, j such that "0 ≤ i ∧ i < 10 ∧ (2i+1=j ∨ read(a,i)=0) ∧ f(read(write(a,i,3), j-2)) ≠ f(j-i+1)"?



# Some Applications of SAI/SMT

- \* Software/Hardware verification
  - \* Model checking, Test-case generation, ...
- \* Theorem proving
- \* Puzzles: Sudoku, Numberlink, Nonogram, etc.
- \* Type checking in Liquid Haskell
  - ★ eg: doubles :: [{x : Int | x >= 0}] → [{x : Int | x `mod` 2 = 0}]
- \* Program Synthesis
- \* and more

# Haskell libraries for SMT solving

# Some Haskell packages for SMT



- \* sbv, smtlib2, simple-smt
- \* z3, bindings-yices, yices-easy, yices-painless
- \* SMT solvers written in Haskell:
  - \* toysolver, Smooth
- \* SMT-LIB2 file parser/printer

\* smt-lib, SmtLib

SMT-LIB2 is a standard input/output format for SMT solvers

# SBV: SMT Based Verification in Haskell

\* SMT library developed by Levent Erkok

### \* It provides:

- High-Level DSL for specifying problems in Haskell, and
- Interfaces to multiple SMT solver backends including Z3, CVC4, Yices, Boolector.
- \* You can install simply using stack/cabal
  - \* "stack install sbv" or "cabal install sbv"

sendMoreMoney :: IO SatResult
sendMoreMoney = sat \$ do
 ds@[s,e,n,d,m,o,r,y] <- mapM sInteger ["s", "e", "n", "d", "m", "o", "r", "y"]
let isDigit x = x .>= 0 &&& x .<= 9
 val xs = sum \$ zipWith (\*) (reverse xs) (iterate (\*10) 1)
 send = val [s,e,n,d]
 more = val [m,o,r,e]
 money = val [m,o,n,e,y]
 constrain \$ bAll isDigit ds
 constrain \$ allDifferent ds
 constrain \$ s ./= 0 &&& m ./= 0
 solve [send + more .== money]</pre>

### MONEY

sendMoreMoney :: IO SatResult
sendMoreMoney = sat \$ do

ds@[s,e,n,d,m,o,r,y] <- mapM sInteger ["s", "e", "n", "d", "m", "o", "r", "y"]
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constrain \$ allDifferent ds
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solve [send + more .== money]</pre>

SMT problem is defined using Symbolic monad, and SMT solving is performed by sat :: Symbolic SBool  $\rightarrow$  10 SatResult

sendMoreMoney :: IO SatResult sendMoreMoney = sat \$ do ds@[s,e,n,d,m,o,r,y] <- mapM sInteger ["s", "e", "n", "d", "m", "o", "r", "y"] let isDigit x = x .>= 0 &&& x .<= 9 val xs = sum \$ zipWith (\*) (reverse xs) (iterate (\*10) 1) send = val [s,e,n,d] more = val [m,o,r,e] money = val [m,o,n,e,y] constrain \$ bAll isDigit ds constrain \$ allDifferent ds constrain \$ s ./= 0 &&& m ./= 0 solve [send + more .== money]

slnteger :: String  $\rightarrow$  Symbolic Slnteger creates integer variable

sendMoreMoney :: IO SatResult
sendMoreMoney = sat \$ do
 ds@[s,e,n,d,m,o,r,y] <- mapM sInteger ["s", "e", "n", "d", "m", "o", "r", "y"]
 let isDigit x = x .>= 0 &&& x .<= 9
 val xs = sum \$ zipWith (\*) (reverse xs) (iterate (\*10) 1)
 send = val [s,e,n,d]
 more = val [m,o,r,e]
 money = val [m,o,n,e,y]
 constrain \$ bAll isDigit ds
 constrain \$ allDifferent ds
 constrain \$ s ./= 0 &&& m ./= 0
 solve [send + more .== money]</pre>

Comparison over symbolic values: we have to use slightly difference operators like (.>=), (&&&). Because Haskell's (>=), (&&) returns Bool, but we want SBool.

sendMoreMoney :: IO SatResult
sendMoreMoney = sat \$ do
 ds@[s,e,n,d,m,o,r,y] <- mapM sInteger ["s", "e", "n", "d", "m", "o", "r", "y"]
 let isDigit x = x .>= 0 &&& x .<= 9
 val xs = sum \$ zipWith (\*) (reverse xs) (iterate (\*10) 1)
 send = val [s,e,n,d]
 more = val [m,o,r,e]
 money = val [m,o,n,e,y]
 constrain \$ bAll isDigit ds
 constrain \$ allDifferent ds
 constrain \$ s ./= 0 &&& m ./= 0
 solve [send + more .== money]</pre>

val ::  $CSIntegerI \rightarrow SInteger$  is defined as in normal Haskell. Thanks to the Num type class.

sendMoreMoney :: IO SatResult
sendMoreMoney = sat \$ do
 ds@[s,e,n,d,m,o,r,y] <- mapM sInteger ["s", "e", "n", "d", "m", "o", "r", "y"]
 let isDigit x = x .>= 0 &&& x .<= 9
 val xs = sum \$ zipWith (\*) (reverse xs) (iterate (\*10) 1)
 send = val [s,e,n,d]
 more = val [m,o,r,e]
 money = val [m,o,n,e,y]
 constrain \$ bAll isDigit ds
 constrain \$ allDifferent ds
 constrain \$ s ./= 0 &&& m ./= 0
 solve [send + more .== money]</pre>

### Actual constraints specification

sendMoreMoney :: IO SatResult sendMoreMoney = sat \$ do ds@[s,e,n,d,m,o,r,y] <- mapM sInteger ["s", "e", "n", "d", "m", "o", "r", "y"] let isDigit x = x .>= 0 &&& x .<= 9 val xs = sum \$ zipWith (\*) (reverse xs) (iterate (\*10) 1) send = val [s,e,n,d] Satisfiable. Model: more = val [m,o,r,e] money = val [m,o,n,e,y] s = 9 :: Integer constrain \$ bAll isDigit ds constrain \$ allDifferent ds e = 5 :: Integer constrain \$ s ./= 0 &&& m ./= 0 n = 6 :: Integer solve [send + more .== money] d = 7 :: Integerm = 1 :: Integer

o = 0 :: Integer

r = 8 :: Integer

y = 2 :: Integer

You need SMT solver Z3 to run the code.

sendMoreMoney :: IO AllSatResult
sendMoreMoney = allSat \$ do
 ds@[s,e,n,d,m,o,r,y] <- mapM sInteger ["s", "e", "n", "d", "m", "o", "r", "y"]
 let isDigit x = x .>= 0 &&& x .<= 9
 val xs = sum \$ zipWith (\*) (reverse xs) (iterate (\*10) 1)
 send = val [s,e,n,d]
 more = val [m,o,r,e]
 money = val [m,o,n,e,y]
 constrain \$ bAll isDigit ds
 constrain \$ allDifferent ds
 constrain \$ s ./= 0 &&& m ./= 0
 solve [send + more .== money]</pre>

By changing sat :: Symbolic SBool  $\rightarrow$  10 SatResult with allSat :: Symbolic SBool  $\rightarrow$  10 AllSatResult

# SBV Summary

# \* This is only one example and sbv includes variety of examples. You should try!

#### Examples

#### BitPrecise

Data.SBV.Examples.BitPrecise.BitTricks Data.SBV.Examples.BitPrecise.Legato Data.SBV.Examples.BitPrecise.MergeSort Data.SBV.Examples.BitPrecise.MultMask Data.SBV.Examples.BitPrecise.PrefixSum

#### CodeGeneration

Data.SBV.Examples.CodeGeneration.AddSub Data.SBV.Examples.CodeGeneration.CRC\_USB5 Data.SBV.Examples.CodeGeneration.Fibonacci Data.SBV.Examples.CodeGeneration.GCD Data.SBV.Examples.CodeGeneration.PopulationCount Data.SBV.Examples.CodeGeneration.Uninterpreted

#### Crypto

Data.SBV.Examples.Crypto.AES Data.SBV.Examples.Crypto.RC4

#### Existentials

Data.SBV.Examples.Existentials.CRCPolynomial Data.SBV.Examples.Existentials.Diophantine

#### Misc

Data.SBV.Examples.Misc.Auxiliary Data.SBV.Examples.Misc.Enumerate Data.SBV.Examples.Misc.Floating Data.SBV.Examples.Misc.ModelExtract Data.SBV.Examples.Misc.NoDiv0 Data.SBV.Examples.Misc.Word4

#### Polynomials

Data.SBV.Examples.Polynomials.Polynomials Puzzles

Data.SBV.Examples.Puzzles.Birthday Data.SBV.Examples.Puzzles.Coins Data.SBV.Examples.Puzzles.Counts Data.SBV.Examples.Puzzles.DogCatMouse Data.SBV.Examples.Puzzles.Euler185 Data.SBV.Examples.Puzzles.Fish Data.SBV.Examples.Puzzles.MagicSquare Data.SBV.Examples.Puzzles.NQueens Data.SBV.Examples.Puzzles.SendMoreMoney Data.SBV.Examples.Puzzles.Sudoku Data.SBV.Examples.Puzzles.Sudoku Data.SBV.Examples.Puzzles.U2Bridge

Data.SBV.Examples.Uninterpreted.AUF Data.SBV.Examples.Uninterpreted.Deduce Data.SBV.Examples.Uninterpreted.Function Data.SBV.Examples.Uninterpreted.Shannon Data.SBV.Examples.Uninterpreted.Sort Data.SBV.Examples.Uninterpreted.UISortAllSat

# toysolver package

- \* I'm implementing some decision procedure in Haskell to leaning the algorithms
  - https://github.com/msakai/toysolver
  - http://hackage.haskell.org/package/toysolver
- \* It contains some algorithms/solvers.
- In particular, it contains a SAT solver 'toysat' and SMT solver 'toysmt'

# Recalling Last Year ...

#### SAT/SMT Y ルバのしくみ

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2015-09-12 Proof Summit 2015

改訂版

### \* At Proof Summit 2015, I talked about how SAT/SMT solver works.

http://www.slideshare.net/sakai/satsmt

 At that time, I already had implemented SAT solver 'toysat', but not implemented SMT solver yet.

\* It triggered my motivation to implement a SMT solver, I worked hard, and finally I did it!

toysat / toysmt

### \* Written in pure Haskell

- \* but implemented in very imperative way
- \* toysat is modestly fast.
  - It was once the fastest among SAT solvers written in Haskell. But now mios by Shoji Narazaki is faster.
- \* toysmt is slow, and has very limited features.

## toysmt

### \* toysat based SMT solver

- \* implementation is really native and notefficient at all
- \* Theories
  - \* Equality and Uninterpreted functions  $\checkmark$
  - \* Linear Real Arithmetic 🗸
  - \* Bit-vector (currently implementing)
  - \* Linear Integer Arithmetic, Array, etc. (not yet)

# toysmt: demonstration

(set-option :produce-models true) (set-logic QF\_UFLRA) (declare-sort U 0) (declare-fun x () Real) (declare-fun f (U) Real) (declare-fun P(U) Bool) (declare-fung(U)U) (declare-fun c () U) (declare-fund()U) (assert (= (P c) (= (g c) c))) (assert (ite (P c) (> x (f d)) (< x (f d))) (check-sat) (get-model) (exit)

### QF\_UFLRA.smt2

# toysmt: demonstration

\$ toysmt QF\_UFLRA.smt2
success

sat ((define-fun P ((x!1 U)) Bool (ite (= x!1 (as @3 V)) true false)) (define-fun c () U (as @3 U)) (define-fun d () V (as @4 V)) (define-fun f ((x!1 U)) Real (ite (= x!1 (as @4 V)) 0 (/ 5555555 1))) (define-fun g ((x!1 V)) V (ite (= x!1 (as @3 V)) (as @3 V) (as @-1 V))) (define-fun x () Real (/ 1 10)))

## For those who do not read SEXP V = {@-1, @1, ..., @4, ...} x = 1/10 : Real c = @3 : V d = @4 : UP(x) = if x = @3 then true else falsef(x) = if x = @4 then 0 else 55555q(x) = if x = @3 then @3 else @-1

# toysmt in SMT-COMP 2016

### QF\_LRA (Main Track)

http://smtcomp.sourceforge.net/2016/results-QF\_LRA.shtml?v=1467876482

Solver	Sequential performance			Parallel performance				Other information
	Error Score	Correctly Solved Score	avg. CPU time	Errors	Corrects	avg. CPU time	avg. WALL time	Unsolved benchmarks
CVC4	0.000	1601.997	61.989	0.000	1601.997	62.004	62.088	20
MathSat5 <sup>n</sup>	0.000	1574.475	109.915	0.000	1574.475	109.957	109.870	64
OpenSMT2	0.000	1510.710	214.228	0.000	1510.710	214.323	214.198	263
SMT-RAT	0.000	1415.279	344.351	0.000	1415.279	344.538	344.329	433
SMTInterpol	0.000	1571.240	118.790	0.000	1572.061	127.002	112.669	54
Yices2	0.000	1593.356	65.700	0.000	1593.356	65.730	65.654	34
toysmt	0.000	1172.885	582.313	0.000	1172.885	582.500	582.280	811
veriT-dev	0.000	1577.045	95.717	0.000	1577.045	95.759	95.701	55
z3 <u>n</u>	0.000	1547.832	157.261	0.000	1547.832	157.327	157.232	145

'toysmt' ended up dead last. But without wrong results! (Thanks to QuickCheck!)

# toysmt: Future work

- \* Fill the gap with state-of-the-art solvers (even a little)
  - \* There're lots of rooms for performance improvement.
  - \* More theories: Bit-vectors, Integer arithmetic, Array, ...
  - \* More features: e.g. Proof-generation
- \* Using 'toysmt' as a backend of 'sbv'.
- \* Re-challenge in next year's SMT-COMP competition.



- \* SAT solvers are amazingly fast for solving many combinatorial problems
- \* SMT is an extension of SAT to handle high-level constraints using specialized solvers.
- \* sbv is a neat Haskell library for using SMT solvers
- \* toysmt is a SMT solver written in Haskell

# Further readings

### How a CDCL SAT Solver works



Masahiro Sakai Twitter: @masahiro\_sakai

#### SAT/SMTYルバのしくみ

酒井 政裕

2015-09-12 Proof Summit 2015

改訂版

#### http://www.slideshare.net/sakai/ how-a-cdcl-sat-solver-works

http://www.slideshare.net/sakai/satsmt

# Further readings

### \* Handbook of Satisfiability

\* A. Biere, M. Heule, H. Van Maaren, and T. Walsh, Eds.

\* IOS Press, Feb. 2009.

 It is a very good book covering variety of topics related to SAT/ SMT.

